

The Logic of Common Ignorance

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Introduction

Quote “Knowledge is a big subject. Ignorance is bigger. . . and it is more interesting.”¹

Claim Ignorance has some surprising properties.

Example Common ignorance.


¹Stuart Firestein, Interview about S. Firestein, *Ignorance: How It Drives Science*, OUP 2012.

Question

- ▶ “Obama calls Trump ignorant about foreign affairs” (Google, August 16, 2016, 8 results).
- ▶ “Trump calls Obama ignorant about foreign affairs” (Google, August 16, 2016, about 135 results).
- ▶ Suppose that at least one of them were right. (Of course, both could be right.)
- ▶ Would this give the group of all humans *common ignorance* about foreign affairs?

Knowing that

To answer this question, we extend the (propositional) logic of individual, shared and common knowledge that A , $TEC_{(m)}$, with a few uncontroversial definitions. $TEC_{(m)}$ applies to a group having members $1, \dots, m$. $TEC_{(m)}$ is well-known and is axiomatized as follows.²

²J.-J. Ch. Meyer and W. van der Hoek, *Epistemic Logic for Computer Science and Artificial Intelligence* (Cambridge: Cambridge University Press, 1995), Ch. 2.1. 

Symbols

- ▶ Individual knowledge that A : $K_i A$, where $1 \leq i \leq m$. $K_i A$ is read as “ i individually knows that A ” or as “ i has individual knowledge that A .”
- ▶ Shared knowledge that A : EA . EA is read as “everyone knows that A ” or as “the group has shared knowledge that A .”
- ▶ Common knowledge that A : CA . CA is read as “it is commonly known that A ” or as “the group has common knowledge that A .”

Axioms and derivation rules

A1 All instances of propositional tautologies.

A2 $\mathbf{K}_i(A \rightarrow B) \rightarrow (\mathbf{K}_iA \rightarrow \mathbf{K}_iB)$.

A3 $\mathbf{K}_iA \rightarrow A$.

A4 $\mathbf{E}A \leftrightarrow \bigwedge_{i=1}^m \mathbf{K}_iA$.

A5 $\mathbf{C}A \rightarrow A$.

A6 $\mathbf{C}A \rightarrow \mathbf{E}CA$.

A7 $\mathbf{C}(A \rightarrow B) \rightarrow (\mathbf{C}A \rightarrow \mathbf{C}B)$.

A8 $\mathbf{C}(A \rightarrow \mathbf{E}A) \rightarrow (A \rightarrow \mathbf{C}A)$.

R1 From A and $A \rightarrow B$ infer B .

R2 From A infer \mathbf{K}_iA .

R3 From A infer $\mathbf{C}A$.

Theorems

- 1.1 $CA \rightarrow EA$ (common knowledge that A implies shared knowledge that A).
- 1.2 $EA \rightarrow K_i A$ (shared knowledge that A implies individual knowledge that A).
- 1.3 $CA \rightarrow K_i A$ (common knowledge that A implies individual knowledge that A).
- †1.4 $K_i A \rightarrow CA$ (individual knowledge that A implies common knowledge that A) is *invalid* [proof: by the semantics].

Intuitively, $CA = \bigwedge_{i \geq 0} E^i A$ (common knowledge that A is the conjunction of A , shared knowledge that A , shared knowledge that the group has shared knowledge that A , and so on).

Knowledge whether/about

Symbols:³

- ▶ Individual knowledge about A : $\Delta_i A = K_i A \vee K_i \neg A$. $\Delta_i A$ is read as “ i individually knows whether A ” or as “ i has individual knowledge about A .”
- ▶ Common knowledge about A : $C_{\Delta} A = C A \vee C \neg A$. $C_{\Delta} A$ is read as “the group has common knowledge about A .”

³See J. Fan, Y. Wang and H. van Ditmarsch, “Contingency and Knowing Whether,” *The Review of Symbolic Logic*, 8:75–107, 2015.

Theorems

2.1 $\mathbf{C}_\Delta A \rightarrow \Delta_i A$ [$(\mathbf{C}A \vee \mathbf{C}\neg A) \rightarrow (\mathbf{K}_i A \vee \mathbf{K}_i \neg A)$] (common knowledge about A implies individual knowledge about A)
[from $\mathbf{C}A \rightarrow \mathbf{K}_i A$ (1.3) by propositional calculus].

†2.2 $\Delta_i A \rightarrow \mathbf{C}_\Delta A$ (individual knowledge about A implies common knowledge about A) is *invalid* [proof: by the semantics].

Ignorance whether/about

Symbols:⁴

- ▶ Individual ignorance about A :

$\nabla_i A = \neg \Delta_i A = \neg K_i A \wedge \neg K_i \neg A$ (individual ignorance about A is the negation of individual knowledge about A). $\nabla_i A$ is read as “ i does not individually know whether A ”, as “ i individually ignores whether A ” or as “ i has individual ignorance about A .”

- ▶ Common ignorance about A :

$C_{\nabla} A = \neg C_{\Delta} A = \neg C A \wedge \neg C \neg A$ (common ignorance about A is the negation of common knowledge about A). $C_{\nabla} A$ is read as “the group has common ignorance about A .”

⁴See Fan, Wang and Van Ditmarsch, “Contingency and Knowing Whether,”
op. cit.

Theorems

3.1 $\nabla_i A \rightarrow \mathbf{C}_{\nabla} A$ [$\neg \Delta_i A \rightarrow \neg \mathbf{C}_{\Delta} A$] (individual ignorance about A implies common ignorance about A) [from $\mathbf{C}_{\Delta} A \rightarrow \Delta_i A$ (2.1) by contraposition].

†3.2 $\mathbf{C}_{\nabla} A \rightarrow \nabla_i A$ (common ignorance about A implies individual ignorance about A) is *invalid* [proof: by the semantics].

Individual ignorance about A is therefore stronger than common ignorance about A . If agents have individual ignorance about A , all groups to which they belong have common ignorance about A .

Answer to question

- ▶ Obama and Trump called each other ignorant about foreign affairs.
- ▶ Suppose that at least one of them were right.
- ▶ Question: would this give the group of all humans *common ignorance* about foreign affairs?
- ▶ Answer: yes, it would, by theorem $\nabla_i A \rightarrow \mathbf{C}_{\nabla} A$ (3.1).

Common ignorance about common ignorance

- ▶ $S5EC_{(m)}$ is $TEC_{(m)}$ plus $\neg K_i A \rightarrow K_i \neg K_i A$ (“ i does not know that A ” implies “ i knows that i does not know that A ”).
- ▶ $S5EC_{(m)}$ has the following theorem.⁵
 - 4.1 $\neg C_{\nabla} C_{\nabla} A$ (there is no common ignorance about common ignorance about A).
- ▶ $TEC_{(m)}$ does not have this theorem, as the semantics shows.
- ▶ The Obama/Trump case seems to show that 4.1 is false.
- ▶ We do have common ignorance about our common ignorance about foreign affairs.
- ▶ $TEC_{(m)}$ is therefore preferable to $S5EC_{(m)}$.

⁵H. Montgomery and R. Routley, “Contingency and Non-Contingency Bases for Normal Modal Logics,” *Logique et Analyse*, 9:318–328, 1966.